

CS370



Symbolic Programming Declarative Programming

LECTURE 3: Syntax and Meaning of Prolog Programs

Jong C. Park

park@cs.kaist.ac.kr

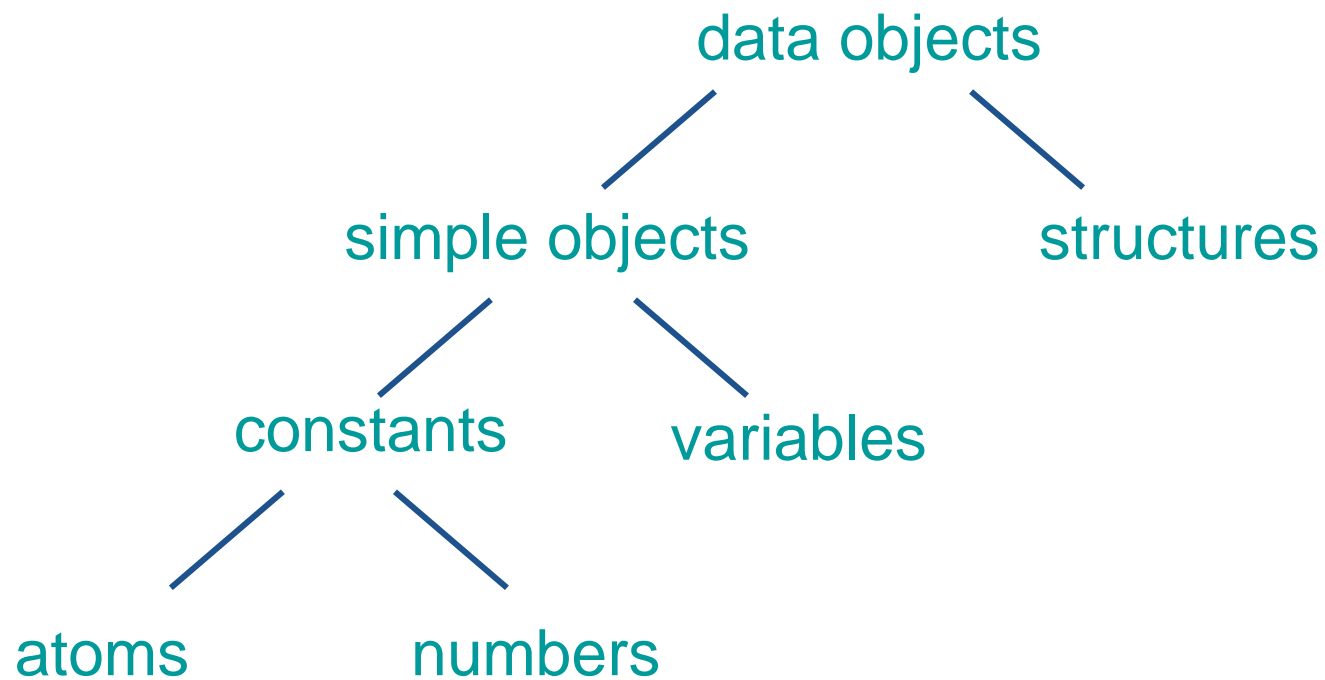
Computer Science Department
Korea Advanced Institute of Science and Technology

<http://nlp.kaist.ac.kr/~cs370>

Syntax and Meaning of Prolog Programs

- ⊙ **Data objects**
- ⊙ **Matching**
- ⊙ **Declarative meaning of Prolog programs**
- ⊙ Procedural meaning
- ⊙ Example: monkey and banana
- ⊙ Order of clauses and goals
- ⊙ The relation between Prolog and logic

Data Objects



Data objects

◎ Atoms

- ◆ Strings of letters, digits and '_', starting with a lowercase letter
 - `anna`, `nil`, `x25`, `x_25`,
`miss_Jones`, `sarah_kerrighan`
- ◆ Strings of special characters
 - `<---->`, `=====>`, `...`, `...`, `::=`
- ◆ Strings of characters enclosed in single quotes
 - `'Tom'`, `'South_America'`, `'Sarah Kerrighan'`

Data objects

© Numbers

◆ Integers

- 1
- 1313
- 0
- -97

◆ Real numbers

- 3.14
- -0.0035
- 100.2

Data Objects

◎ Variables

- ◆ Strings of letters, digits, and '_' that start with an uppercase letter or '_'
 - `X`, `Result`, `Object2`, `_x23`, `_23`
- ◆ Anonymous variable: `_`
 - `hasachild(X) :- parent(X,_)`
 - `somebody_has_child :- parent(_,_).`
- ◆ Lexical scope of a variable

Data Objects

◎ Structures

- ◆ Structures with 3 components
 - Examples
 - `date(1,may,2001)`
 - `date(Day,may,2001)`
 - ◆ Terminology
 - Term
 - Functor
 - Argument
 - Principal functor

`father(bob)`

Data Objects

© Structures

◆ Geometric objects in 2D space

▪ Examples

- $P1 = \text{point}(1,1)$
- $P2 = \text{point}(2,3)$
- $S = \text{seg}(P1, P2)$
 $= \text{seg}(\text{point}(1,1), \text{point}(2,3))$
- $T = \text{triangle}(\text{point}(4,2), \text{point}(6,4), \text{point}(7,1))$

▪ Any other ways?

-

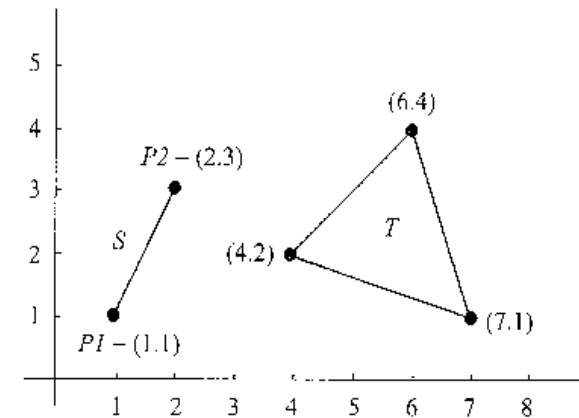


Figure 2.3 Some simple geometric objects.

Data Objects

© Structures

- ◆ Representing objects in 3D space
 - Alternatives
 - `point3(X,Y,Z)`
 - `point(X,Y,Z)`
 - What are the pros and cons?
 -

Matching

◎ Two terms match if

- ◆ they are identical, or
- ◆ the variables in both terms can be instantiated to objects in such a way that after the substitution of variables by these objects the terms become identical

◎ Examples

- ◆ `date(D,M,2001)` and `date(D1,may,Y1)`?
- ◆ `date(D,M,2001)` and `date(D1,M1,1444)`?
- ◆ `date(X,Y,Z)` and `point(X,Y,Z)`?

Matching

⊙ Given two terms S and T :

- ◆ If S and T are $\text{atom}(x)$ then S and T match only if they are the same object.
- ◆ If S is a $\text{var}(x)$ and T is anything, then they match, and S is instantiated to T . Conversely, if T is a variable then T is instantiated to S .
- ◆ If S and T are $\text{functor}(f, \text{args})$ then they match only if they have the same principal functor and all their corresponding components match.

Matching

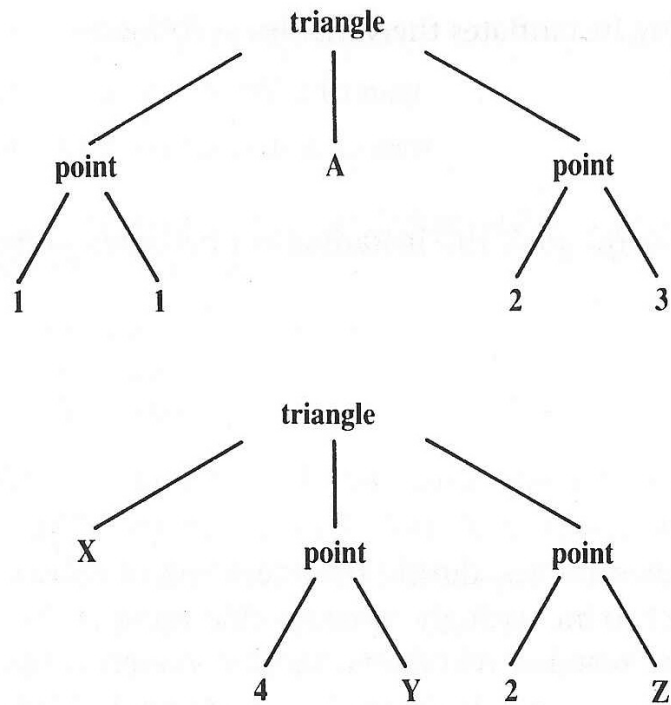


Figure 2.7 Matching $\text{triangle}(\text{point}(1,1), A, \text{point}(2,3)) = \text{triangle}(X, \text{point}(4,Y), \text{point}(2,Z))$.

Matching

◎ Example use of matching

- ◆ Recall the representation for line segments
 - $S = \text{seg}(P1, P2) = \text{seg}(\text{point}(1, 1), \text{point}(2, 3))$
- ◆ Define a piece of program for recognizing horizontal and vertical line segments
 - `vertical(seg(point(X, Y), point(X, Y1))).`
 - `horizontal(seg(point(X, Y), point(X1, Y))).`

?-

Prolog programs

◎ Meanings

- ◆ $P :- Q, R.$
- ◆ declarative meaning
 - P is true if Q and R are true.
 - From Q and R follows P.
- ◆ procedural meaning
 - To solve problem P, first solve the subproblem Q and then the subproblem R.
 - To satisfy P, first satisfy Q and then R.

Prolog programs

◎ Instances and variants of a clause

- ◆ Example
 - `hasachild(X) :- parent(X,Y).`
- ◆ A **variant** of a clause C is the clause C where each variable is substituted by another variable.
 - `hasachild(A) :- parent(A,B).`
 - `hasachild(X1) :- parent(X1,X2).`
- ◆ An **instance** of a clause C is the clause C with each of its variables substituted by some term.
 - `hasachild(peter) :- parent(peter,Z).`

Prolog programs

⊙ **Given a program and a goal G , the declarative meaning says:**

- ◆ A goal G is true iff:
 - there is a clause C in the program such that
 - there is a clause instance I of C such that
 - the head of I is identical to G , and
 - all the goals in the body of I are true.

⊙ **A **question** is true if all of its goals are true for the same instantiation of variables.**

Prolog programs

◎ Conjunction and disjunction of goals

- ◆ Conjunction of goals

- $P :- Q, R.$

- ◆ Disjunction of goals

- $P :- Q; R.$

- $P :- Q.$

- $P :- R.$

- $P :- Q, R; S, T, U.$

- $P :- Q, R.$

- $P :- S, T, U.$

Summary

- ⊙ Simple objects in Prolog are **atoms**, **variables** and **numbers**.
- ⊙ Structured objects (**structures**) are used to represent objects that have several components.
- ⊙ Structures are constructed by means of **functors**. Each functor is defined by its **name** and **arity**.
- ⊙ The **type** of object is recognized entirely by its syntactic form.

Summary

- ⊙ The **lexical scope** of variables is one clause. Thus the same variable name in two clauses means two different variables. **Structures** can be naturally pictured as **trees**. Prolog can be viewed as language for processing trees.
- ⊙ The **matching** operation takes two terms and tries to make them identical by instantiating the variables in both terms.
- ⊙ Matching, if it succeeds, results in the most general instantiation of variables.

Summary

- ⊙ The **declarative semantics** of Prolog defines whether a goal is true wrt a given program, and if it is true, for what instantiation of variables it is true.
- ⊙ A **comma** between goals means the conjunction of goals.
- ⊙ A **semicolon** between goals means the disjunction of goals.